

rows of A . A matrix in $\mathbb{C}^{n \times m}$ is of *full rank* when its rank is equal to the smallest of m and n . A fundamental result of linear algebra is stated by the following relation:

$$\mathbb{C}^n = \text{Ran}(A) \oplus \text{Ker}(A^T). \quad (1.17)$$

The same result applied to the transpose of A yields $\mathbb{C}^m = \text{Ran}(A^T) \oplus \text{Ker}(A)$.

A subspace S is said to be *invariant* under a (square) matrix A whenever $AS \subset S$. In particular, for any eigenvalue λ of A the subspace $\text{Ker}(A - \lambda I)$ is invariant under A . The subspace $\text{Ker}(A - \lambda I)$ is called the *eigenspace* associated with λ and consists of all the eigenvectors of A associated with λ , in addition to the zero vector.

1.7 Orthogonal Vectors and Subspaces

A set of vectors $G = \{a_1, a_2, \dots, a_r\}$ is said to be *orthogonal* if

$$(a_i, a_j) = 0 \quad \text{when } i \neq j.$$

It is *orthonormal* if, in addition, every vector of G has a 2-norm equal to unity. A vector that is orthogonal to all the vectors of a subspace S is said to be *orthogonal* to this subspace. The set of all the vectors that are orthogonal to S is a vector subspace called the *orthogonal complement* of S and denoted by S^\perp . The space \mathbb{C}^n is the direct sum of S and its orthogonal complement. Thus, any vector x can be written in a unique fashion as the sum of a vector in S and a vector in S^\perp . The operator that maps x into its component in the subspace S is the *orthogonal projector* onto S .

Every subspace admits an orthonormal basis that is obtained by taking any basis and *orthonormalizing* it. The orthonormalization can be achieved by an algorithm known as the *Gram-Schmidt process*, which we now describe.

Given a set of linearly independent vectors $\{x_1, x_2, \dots, x_r\}$, first normalize the vector x_1 , which means divide it by its 2-norm, to obtain the scaled vector q_1 of norm unity. Then x_2 is *orthogonalized* against the vector q_1 by subtracting from x_2 a multiple of q_1 to make the resulting vector orthogonal to q_1 ; i.e.,

$$x_2 \leftarrow x_2 - (x_2, q_1)q_1.$$

The resulting vector is again normalized to yield the second vector q_2 . The i th step of the Gram-Schmidt process consists of orthogonalizing the vector x_i against all previous vectors q_j .

ALGORITHM 1.1. Gram-Schmidt

1. Compute $r_{11} := \|x_1\|_2$. If $r_{11} = 0$ Stop, else compute $q_1 := x_1/r_{11}$
2. For $j = 2, \dots, r$, Do
3. Compute $r_{ij} := (x_j, q_i)$ for $i = 1, 2, \dots, j-1$
4. $\hat{q} := x_j - \sum_{i=1}^{j-1} r_{ij}q_i$
5. $r_{jj} := \|\hat{q}\|_2$
6. If $r_{jj} = 0$ then Stop, else $q_j := \hat{q}/r_{jj}$
7. EndDo

It is easy to prove that the above algorithm will not break down; i.e., all r steps will be completed iff the set of vectors x_1, x_2, \dots, x_r is linearly independent. From lines 4 and 5, it is clear that at every step of the algorithm the following relation holds:

$$x_j = \sum_{i=1}^j r_{ij} q_i.$$

If $X = [x_1, x_2, \dots, x_r]$, $Q = [q_1, q_2, \dots, q_r]$, and R denotes the $r \times r$ upper triangular matrix whose nonzero elements are the r_{ij} 's defined in the algorithm, then the above relation can be written as

$$X = QR. \quad (1.18)$$

This is called the QR decomposition of the $n \times r$ matrix X . From what was said above, the QR decomposition of a matrix exists whenever the column vectors of X form a linearly independent set of vectors.

The above algorithm is the standard Gram-Schmidt process. There are alternative formulations of the algorithm that have better numerical properties. The best known of these is the modified Gram-Schmidt (MGS) algorithm.

ALGORITHM 1.2. MGS

1. Define $r_{11} := \|x_1\|_2$. If $r_{11} = 0$ Stop, else $q_1 := x_1/r_{11}$
2. For $j = 2, \dots, r$, Do
3. Define $\hat{q} := x_j$
4. For $i = 1, \dots, j-1$, Do
5. $r_{ij} := (\hat{q}, q_i)$
6. $\hat{q} := \hat{q} - r_{ij}q_i$
7. EndDo
8. Compute $r_{jj} := \|\hat{q}\|_2$
9. If $r_{jj} = 0$ then Stop, else $q_j := \hat{q}/r_{jj}$
10. EndDo

Yet another alternative for orthogonalizing a sequence of vectors is the Householder algorithm. This technique uses Householder reflectors, i.e., matrices of the form

$$P = I - 2ww^T, \quad (1.19)$$

in which w is a vector of 2-norm unity. Geometrically, the vector Px represents a mirror image of x with respect to the hyperplane $\text{span}\{w\}^\perp$.

To describe the Householder orthogonalization process, the problem can be formulated as that of finding a QR factorization of a given $n \times m$ matrix X . For any vector x , the vector w for the Householder transformation (1.19) is selected in such a way that

$$Px = \alpha e_1,$$

where α is a scalar. Writing $(I - 2ww^T)x = \alpha e_1$ yields

$$2w^T x w = x - \alpha e_1. \quad (1.20)$$

المخافة الباردة :
 "تنتج الصيغة المتأثرة"
 هواريزميو غرام - شينيت

1- compute $r_{11} := \|x_1\|_2$

إذا كانه :

$$r_{11} = 0$$

لتوقف

بالا :

compute $q_1 := x_1 / r_{11}$

2- for $j=2, \dots, n$ do

3- compute $r_{ij} := (x_j, q_i)$ for $i=1, 2, \dots, j-1$

4- $\hat{q} = x_j - \sum_{i=1}^{j-1} r_{ij} q_i$

5- $r_{jj} := \|\hat{q}\|_2$

6- IF $r_{jj} = 0$ then stop, else $q_j := \hat{q} / r_{jj}$

I- End DO

"الصيغة الكاملة محسنة"

من أجل ان نثبت ان الخوارزمية السابقة لن تقف

تدريج كل المحصولات ، سيقف اذا وفقط اذا كانت المتجهات x_1, x_2, \dots, x_r مستقلة خطياً

من الصيغ 4 و 5 ، دافع انه في كل خطوة من الخوارزمية تتوقف العلاقة التالية

$$x_j = \sum_{i=1}^j r_{ij} q_i$$

اذا كانت $X = [x_1, x_2, \dots, x_r]$ ، $Q_n = [q_1, q_2, \dots, q_r]$ و كانت R مصفوفة

مُثلثة عليا من الرتبة $r \times r$ والتي عندها النصف السفلي هو عبارة عن r_{ij} المعرفة في الخوارزمية. عندها العلاقة السابقة يمكن كتابتها بالشكل

$$X = QR \quad (1.18)$$

إن هذا يعبر بالتحال QR للمصفوفة X من المرتبة $n \times r$.
 نستطيع ما فعل سابقاً، إن QR للمصفوفة X يوجد وفقاً لتلك المجموعة
 المصفوفة X مستقلة خطياً من الأعمدة.
 إن الخوارزمية السابقة هي طريقة غرام-شmidt المتكررة.
 إن هناك مميزات بديلة للخوارزمية والتي تمكننا من إيجاد عددية أفضل
 أكثرها هي طريقة غرام-شmidt المعدلة.

خوارزمية ALGORITHM 1.2. MGS

هناك ثلاث خطوات بدلاً من التقسيم متتالية من الأعمدة في خوارزمية جاورس
 هذه العملية تستخدم عملاً n لها n عمود
 تدعى w ، z n من n كما:

$$P = I - zw^T \quad (1.19)$$

حيث w هو متجه ذو نقيص أقل من n واحد.

هناك شيئاً السامع PX يُمثل الصورة المباشرة لـ X المتعلقة بالعضاء المعامد
 للعضاء المولد بـ w .

لدينا عملية تسمى جاورس-هولدر، يمكن أن نضع المسألة إلى إيجاد QR
 للمصفوفة المربعة X من المرتبة $n \times m$
 من أجل أي متجه x ، السامع w

$$PX = \alpha e_1$$

where α is a scalar. writing $(I - zw^T)x = \alpha e_1$ yields

$$zw^T x = x - \alpha e_1$$

- السنته المباشرة السادة -